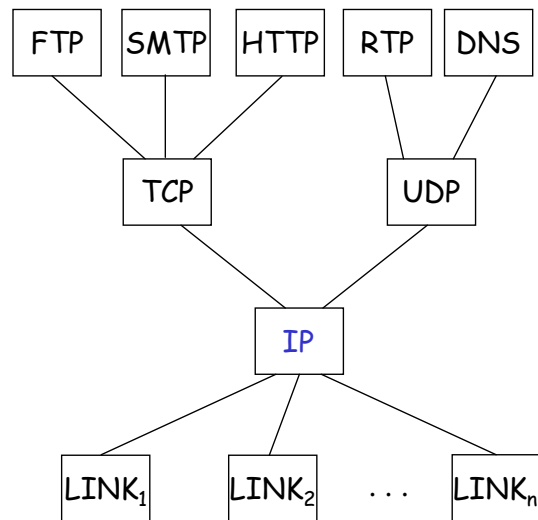


Back to the Future Part 4: The Internet

Simon S. Lam
Department of Computer Sciences
The University of Texas at Austin

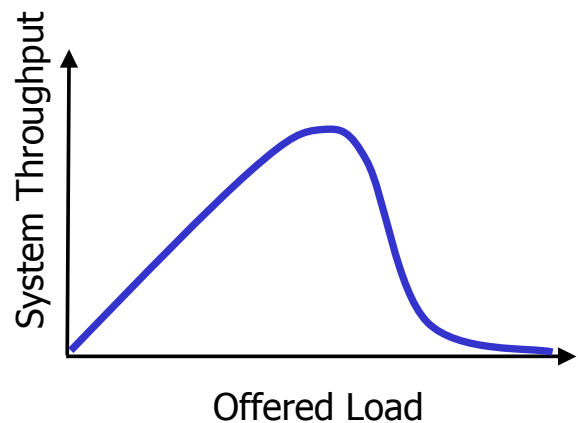
IP won the networking race

- Many competitors in the past
 - SNA, DECnet, XNS
 - X.25, ATM, CLNP
- IP provides end-to-end delivery of datagrams, **best-effort service** only
- IP can use **any link-layer technology** that delivers datagrams (packets)



IP's underlying model is a network of queues

- Revolutionary change from the **circuit** model (Kleinrock 1961)
 - Each packet is routed independently using its destination IP address
 - No concept of a flow between source and destination, no flow state in routers
- Unreliable channels, limited buffer capacity
- Prone to **congestion collapse**



Congestion collapse—ALOHA channel

□ The *ALOHA System* (Abramson 1970)

□ Poisson process assumption

○ pure ALOHA throughput

$$S = G e^{-2G} \quad (\text{Abramson})$$

○ slotted ALOHA throughput

$$S = G e^{-G} \quad (\text{Roberts})$$

□ *ARPANET Satellite System* (1972)

Congestion collapse—ALOHA channel (cont.)

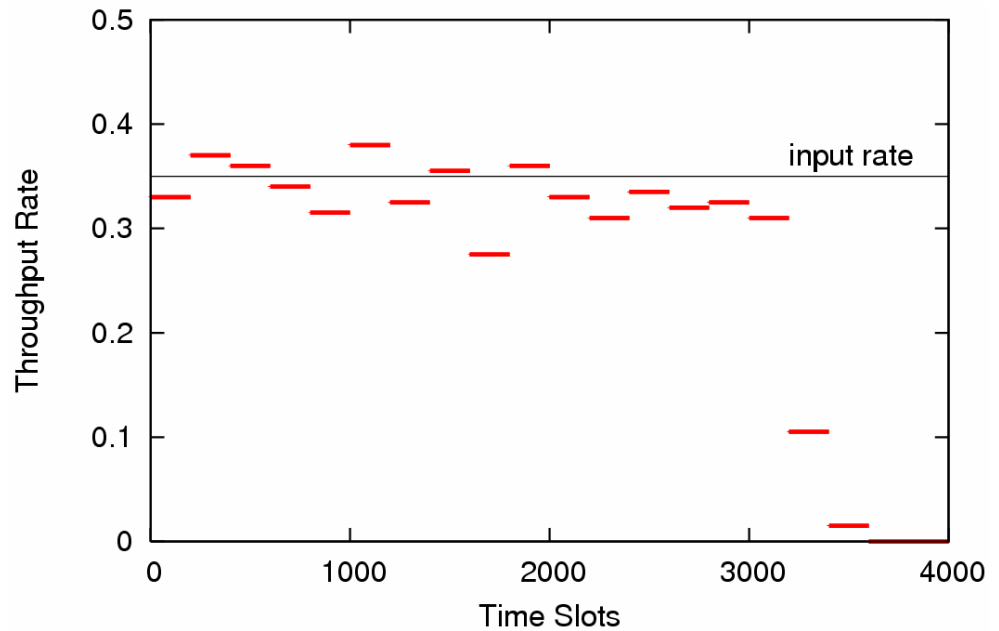
- The early ALOHA System was vastly **under-utilized**

- **Backoff** algorithm for slotted ALOHA (Kleinrock-Lam 1973)
 - Retransmit a collided packet randomly into one of **K** future time slots
 - Poisson process assumption implies
 $K \rightarrow \infty$

Congestion collapse—ALOHA channel (cont.)

□ ASS Note 48 (Lam 1973)

K = 15

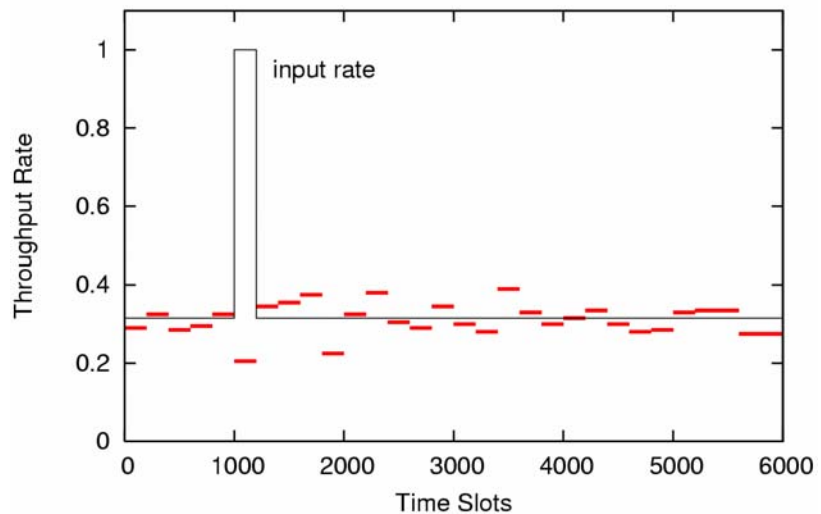


Congestion collapse—ALOHA channel (cont.)

- **Adaptive backoff** algorithm (Lam 1974):
Retransmit a packet with m previous collisions into $K(m)$ slots, where $K(m)$ is monotonically nondecreasing in m

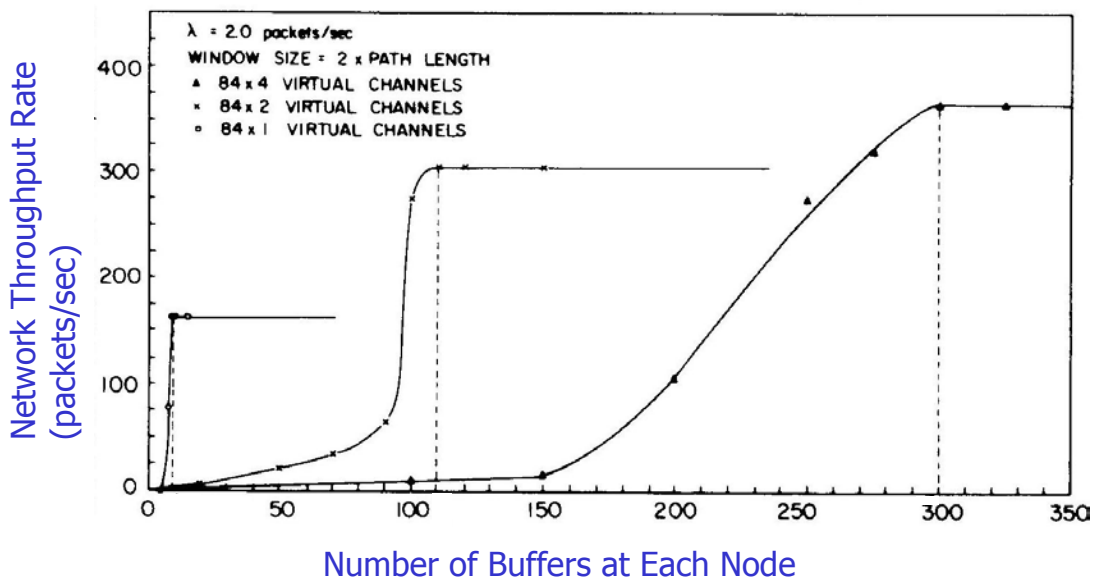
$$K(1) = 10$$

$$K(m) = 150, m \geq 2$$



Congestion collapse—packet networks

- Network of queues with limited buffers
 - Static window flow control in TCP not sufficient (Lam, late 1970s)



Internet congestion control

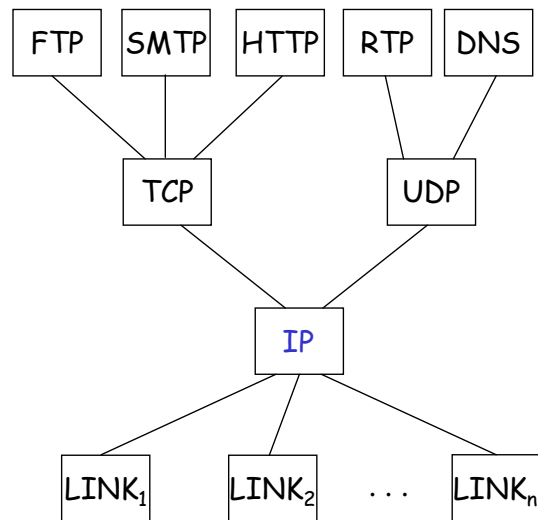
- Van Jacobson's algorithms for TCP congestion control (late 1980s)
 - main reason for stability of the current Internet
- UDP does not perform congestion control
 - preferred by voice and video applications
- More and more voice and video traffic will impact Internet stability
 - I believe in differentiating voice and video flows as well as flow admission control

Efforts to extend/replace IPv4 (past 15 years)

- ❑ IP multicast
- ❑ QoS support - IntServ, RSVP, DiffServ
- ❑ Active Networks research program of DARPA
- ❑ IPsec - retrofitting IP with security
- ❑ IPv6 - replacing IPv4
 - 128 bit IP address
 - flow concept to support QoS
- ❑ Mobile IP
- ❑ . . .

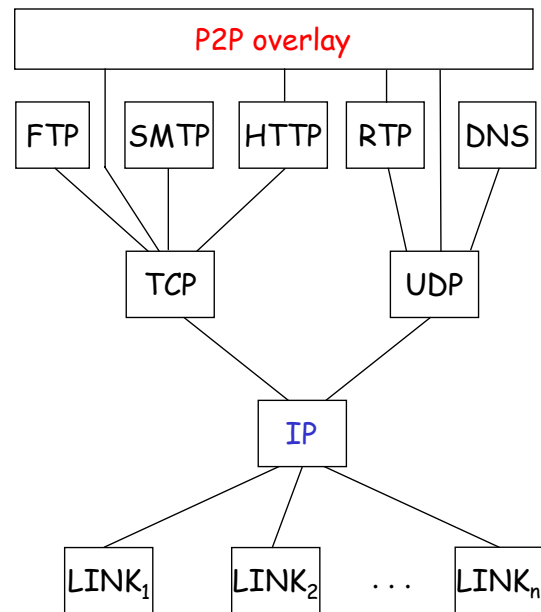
Don't mess with IP?

- In recent years, the research community has moved on to other areas
 - P2P overlays
 - Wireless (ad hoc networks, sensor networks, satellite networks)
 - Measurements
- But the IP foundation, currently relying on over-provisioning, still needs work



Is P2P overlay a panacea?

- ❑ P2P overlay supports many **new** distributed applications
- ❑ P2P overlay is **inefficient** in its use of underlying Internet resources
- ❑ P2P overlay does not directly address **IP's** **foundational issues** (stability, QoS)



IP as the future universal interface

- Payload in the form of IP packets to enable new applications across different telecom networks

analog → digital → packet

- Recent news

- Microsoft unveiled plans for developing IPTV (October 2003)
- SBC to invest \$6 billion on fiber to home for TV services (June 2004)

Changing network traffic mix

- Much more voice and video traffic
- Current traffic of a major telecommunications carrier
 - Circuit switched voice 1.2 petabytes/day
 - Internet traffic 1.5 petabytes/day
- Television services over IP
 - Back of the envelope calculation:
4-8 Mbps, 10,000 seconds/day, 10^8 TV sets
→ 500-1000 petabytes/day to end users

A pragmatic approach

1. Learn from the evolution of Ethernet

Ethernet technology today is very different from Ethernet technology 20 years ago.

- Transmission rates: 10 Mbps, 100 Mbps, 1 Gbps, 10 Gbps
- Switching protocols:
 - CSMA/CD on a cable,
 - CSMA/CD on a hub,
 - collision-free switching,
 - full-duplex point-to-point, both WAN and LAN
- A variety of coding techniques and media

Only the **Ethernet frame interface** remains the same.

A pragmatic approach (cont.)

2. Accept **in some form** a competing idea that has been vanquished again and again by IP but refuses to die and go away:

virtual circuit packet switching

- X.25 1970s - 1990s [L. Roberts]
- Frame Relay 1980s - present
- ATM early 1990s - present
- Label switching, MPLS late 1990s - present

Real circuit switching also under IP

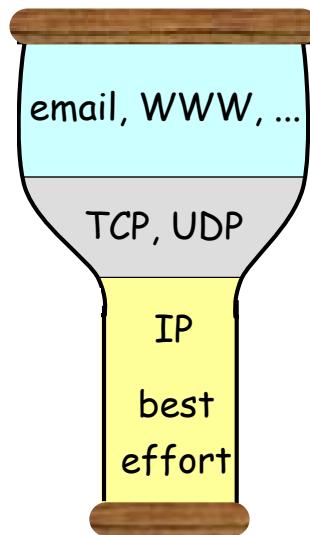
- IP over SONET
- GMPLS forwarding based on TDM time slot, wavelength, or optical port
 - Nesting of label switched paths (LSPs) reminiscent of multiplexing/demultiplexing hierarchy in telephone networks

IP should be a "big tent"

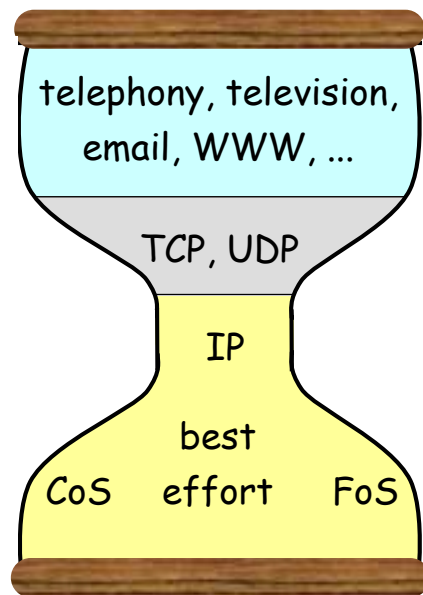
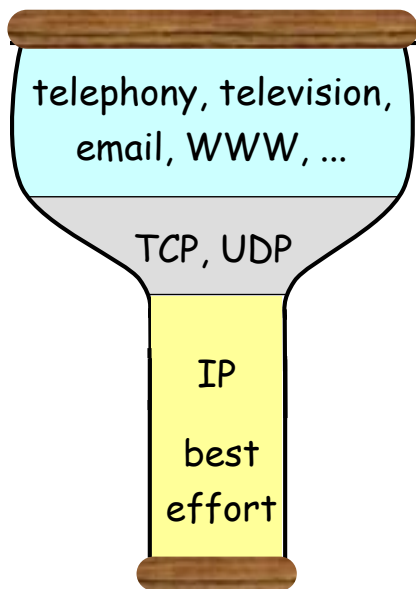
- ❑ To "rule" the world of communications, IP has to attend to the needs of new constituents (voice, video)
 - multiple services to support diverse applications
- ❑ For the research community, **over-provisioning** should be considered a temporary fix, not a permanent solution
- ❑ While the core is over-provisioned, access paths to the core are not

The hourglass shape reconsidered

- Although link-layer technologies are diverse, including virtual and real circuits, only best effort service is available to Internet applications



IP needs a broader base



CoS **Class-oriented Service**

FoS **Flow-oriented Service**

What is flow-oriented service?

- ❑ Dynamic signaling, flow admission control, flow state, quantitative QoS metric
 - Flows subject to admission control rather than random packet drop
 - Make use of virtual or real circuits in link layer

- ❑ Am I reviving IntServ?
Not exactly

- ❑ Per-flow state and dynamic signaling are not scalable!
I know

Two good engineering ideas for voice and video

1. Flow aggregation

- Current examples
Virtual paths in ATM, Label stacks in MPLS, RSVP aggregation
- Each flow is routed along a sequence of "virtual channels" each of which carries a flow aggregate
- Flow aggregation reduces state information and signaling overhead thus improving scalability
 - In the extreme case, a router has just two flow aggregates (voice and video) for each outgoing channel

Two good engineering ideas (cont.)

2. Statistical guarantee

- A natural service guarantee for a voice or video flow is the flow's loss probability for a given packet delay bound

$$\text{Prob}[\text{packet delay} > x] < \varepsilon$$

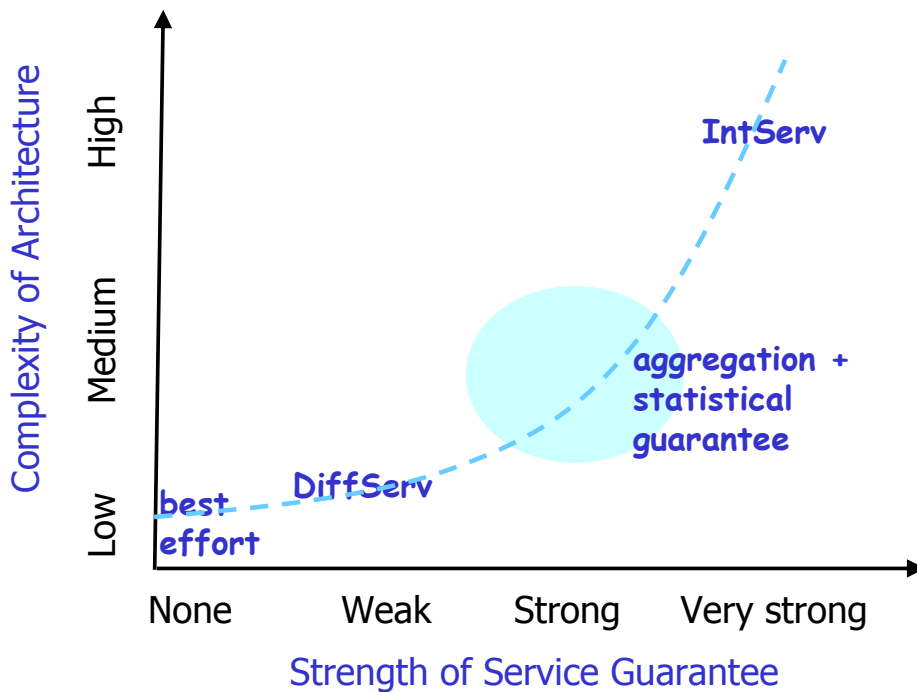
Very hard problem! Substantial research in the past but needs much more work to be applicable.

- **Statistical multiplexing gain** from flow aggregation

Major research issues

- How to derive service guarantee of a flow from the service guarantee provided to a flow aggregate?
- Dynamic configuration and provisioning of virtual channels for flow aggregates.
- How to efficiently compute the end-to-end statistical guarantee to a flow, under practical assumptions?
 - design, modeling, analysis
 - approximation methods
 - measurement-based techniques

QoS research is not done



Inter-provider QoS is a major challenge

- ❑ Business and legal issues
- ❑ Framework for competitive ISPs to cooperate
 - A quantitative QoS metric for inter-provider agreement
 - A small set of standardized traffic specs for voice and video
 - ...
- ❑ It would be nice to have a **de facto standard!**

How to get one?

- The SSL model
 - **Application-driven**—the need to secure web transactions
 - Now SSL used for other applications as well
- The FedEx model
 - **Profit-driven**—someone takes risk

Possible scenario: Some large ISP takes risk and provides QoS services to a large part of the Internet. Success leads to universal global coverage and a de facto standard.

Conclusions

- For the research community, over-provisioning should not be considered a solution
- To become the universal telecom interface, IP needs to be a "big tent"
 - A **flow-oriented service** needed to support television and telephony services
- QoS research is not done
 - Flow aggregation and statistical guarantee merit further investigation

Conclusions (cont.)

□ From history

- Internet—almost 30 years from initial research to commercial deployment
- Packet radio—about 25 years
- QoS research began in late 1980s

Widespread commercial deployment of QoS within 10 years!