

# The Internet Architecture: Its Future and Why it Matters

David Cheriton  
Computer Science Department  
Stanford University

# Internet Architecture

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- What: principles, protocols and structure for highly scalable digital communication
- Principles
  - Application state at endpoints
    - fate-sharing and otherwise soft net. state (D.Clark88)
  - One thin-waist (IP) for simple end-to-end connectivity
    - Multiple types of traffic
  - No off-path components
  - Liberal in what you receive; conservative in what you send

***An amazing accomplishment***

***no thanks to me***

# Internet Architecture: what it provides

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- Properties
  - Survivability: Intermediate nodes can crash and reboot w/o loss of application state
  - Simple to achieve connectivity for different applications
    - Just implement IP plus transport/application protocols
  - Have path, will communicate
  - Interoperability: not need for perfect implementation
- Applications build/rely on these properties

***So, architecture provides properties,  
but only if you are faithful to it***

# The Future

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- Internet-enabled devices are everywhere
- Internet connectivity is ubiquitous
- Internet bandwidth is plentiful
- Special-purpose networks go extinct
  - No separate telephone, TV, SCADA networks
- All critical systems on the public Internet
  - Global financial systems
  - Power distribution systems
  - Air traffic control
  - . . .

***Triumph: unified general-purpose communication  
or is it: a disaster waiting to happen?***

# Air Traffic Control on the Public Internet!

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Crazy? No, because there is no alternative:

- Internet technology: Ethernet switches, IP routers, etc.
  - Market: best, lowest-cost products
  - Staffing: good people know IP, etc.
- Public Internet: Really a collection of ISPs
  - Cost: lowest cost WAN connectivity
  - Availability: expert operators with lots of redundant connectivity
- But how about separation at layer 1?
  - Different colors for Internet ATC (I-ATC)
  - But where does the control plane for the optical run?
    - Single point of failure or public Internet?

***I-ATC is inevitable! And frightening***

# The Internet Architecture: Why it matters?

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The architecture allows us to state properties of the Internet and meet application requirements

- E.g. how to configure to meet I-ATC requirements?

If reality departs from architecture, properties are lost or unknown

- E.g. Ad hoc firewalling and NAT break end-to-end connectivity and reliability

If the architecture is wrong - can fail catastrophically

- The largest, most attractive “asset” to attack in the history of mankind

***It matters too much to be ignored or wrong***

# Unfortunately, it is both

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Ignored? Many violations of the architecture:

- What connectivity can a new wide-area Internet application assume?
  - Port 80 HTTP where the headers “look like” normal HTTP headers, going through NAT
    - Or maybe nothing because of DDoS, route flaps, etc.
  - No end-to-end addressing or reliability
- Dependences on off-path DNS server, root CA

Wrong?

- Current Internet could not work without the above

***A New & Better Internet Architecture is required***

# Trust and Technologies

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- New technologies develop, focused on improving features, performance and cost, however:
- The limit of most technologies is TRUST
- 250 MPH car: can build it, who do you trust to drive?
- Nuclear power plant: most efficient power but limited by trust in who builds and who operates
- GM Foods – we can grow them, will you eat?

**Challenge: Internet architecture trusted to support critical infrastructure systems**



# Internet ATC Requirements

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Very high availability, even under attack:

- Multiple disjoint paths between end-systems with fast fail-over
- Protection against DDoS
- Packet trace-ability – what source
- NOT Performance – low data rate
- NOT Confidentiality – in fact, open to be safe!

Other critical systems have same requirements

***None supported by current architecture;***

***Oh, but ... the work on Internet security!***

# You want security, I have a “solution”

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It's just that it:

- Has a single point of failure
- Is not testable
- Relies on negative acks, not positive acks
- Requires a costly complex implementation that is not understandable by most people
- Does not scale

***Dead-on-arrival in the Internet community?***

***No, it just needs good “packaging”***

# The “Solution”: PKI Certificates

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- Single point of failure
  - Loss of secrecy of private key of root CA
  - Flooding attacks
- Is not testable
  - No way to test if a key is secret
- Uses negative acks, not positive acks
  - Send out nacks in CRLs as part of revocation
- Costly complex implementation
  - PKE, signing, X.509, off-line CAs, CRLs, etc.
- Does not scale: off-line root CA for “security”

***This is Internet security? I don't feel secure!***

# Where did we go wrong

Dictionary: security == safety

- Security was hijacked to mean confidentiality
- Confidentiality was hijacked to mean encryption
  - Same for authentication
- Encryption only “understood” by cryptographers
- So, Internet security delegated to cryptographers
  - Cryptographers are algorithm designers
- Result: Standardized metaprotocols so poor interoperability, no safety, lots of overhead, single point of failure, no useful properties

**Secrecy does *not* scale**

***A secure system needs a system design***

# You want e2e reliability, I have a “solution”

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It's just that it:

- Doesn't provide end-to-end reliability
- Increases exposure to flooding DoS attacks
- Still a design-in-progress after 10 years
- Will take forever to deploy
- Hasn't been evaluated relative to alternatives

***Surely, a non-starter in the Internet community***  
***No, just needs some good marketing, ardent followers and government mandates***

# The “Solution”: IPv6

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- No end-to-end reliability for named endpoints
  - Name-to-address binding can change w/ DHCP
- Exposure to flooding DoS attacks
  - Requires double forwarding/lookup bandwidth
- It is still a design-in-progress after 10 years
  - Addressing architecture, renumbering, mobility, flows
- It will take forever to deploy and makes things worse in the mean time – breaks IP thin waist
  - Upgrading 200 million hosts? IPv4<->IPv6 ?
- No evaluation of alternatives
  - Like change the transport checksum computation?

***An enormous effort in the wrong direction***

# Where did we go wrong?

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- Back in the 1970s - using IP addresses to identify end-system state
  - an IP address identifies an interface on host on particular (sub)network at a particular time
  - IPv6 – further ties it to a particular ISP’s network
  - But state reachable by different interfaces/networks
- Again in the 1990’s, by “believing” e2e IP addresses had some useful semantics

***Reliability requires semantics;***

***IP addresses are transient routing tags,  
nothing more***

# You want routing: I have a “solution”

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It's just that:

- It depends on global trust and competence
- It must be operated at less than 1/1000th of real speed to be stable
- Forces you to upgrade your router as the Internet grows but provides you no benefit
- You have no control beyond first hop (and last I/F)

***Surely, we would never implement . . .  
wrong again!***



# The “Solution”: (secure) BGP

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- global trust and competence
  - Shared world model: believe updates from your peers
  - Signed updates so you can “trust” your peers
- Operated at 1/1000th of real speed for stability
  - 30 second damping to avoid oscillations
- Non-scalable cost
  - Every router stores/recomputes all routes after updates
- You have no control beyond first hop
  - Source routing is disabled by ISPs

***A large open loop dynamic control system  
Defying conventional engineering or ...?***

# Internet WAN Traffic Load

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- Total WWW bandwidth, world-wide
  - P. Danzig 2000 estimate: 250 Gbps!
  - P. Danzig 2003 estimate: 250 Gbps!
  - WWW is half of internet traffic
    - P2P “file sharing” and spam is the rest
- 1/2 single terabit router for entire known universe
- Not an issue even if larger by factor of 10 or more
- Moreover
  - 10 G Ethernet coming down in price
  - lots of dark fiber

***Wide-area bandwidth is not the problem  
wide-area business models are***

# This is all very depressing for I-ATC

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- The Internet architecture is wrong
- The new developments do not address issues
- Research is focused “elsewhere”
- Critical systems will fail with terrible consequences  
when a massive Internet failure happens

***Can we avoid disaster?***

***Let's reboot***

# Cisco: How to sell a router

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- Early days of Cisco: how to get someone to buy a router?
  - Already had connectivity
  - International Ethernets
- Selling point: routers limit Ethernet broadcast storms
  - STP loops, misconfigs would bring down the whole wide-area Ethernet
  - You don't need a router to forward packets
  - You need it to (selectively) not forward packets

***The router as a point of network control***

# Routing as a Broadcast Overlay

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- “Shared world” model of routing – topology info sent everywhere
- Parallel to L2 packet bcast everywhere on unknown address
  - L2 proliferate packet vs. L3 proliferate routing info
  - L2 proliferate packet garbage vs. L3 proliferate routing garbage
- Damage: routing blackhole or grey hole

***The router needs to filter out routing misinformation and select the route, without depending on all other routers***

# Feedback-based Routing

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- Each access router
  - Gets potential routes from “broadcast” topology updates
  - Monitors packet traffic over routes plus sends probes to check potential routes
  - Filters out bad routes, only uses routes known to work
  - Controls packet paths with source routing
- Use feedback, like most engineered dynamic control systems

***Local control and no need for global trust, assuming source routing***

# Source Routing

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- Control the (loose) route each packet takes
- WRAP: Wide-area Relay Addressing Protocol
  - Specifies loose source route for packet
  - Shim protocol over IPv4
- But also, fosters competition among ISPs
- But also, supports NAT Inter-realm addressing
- But also, more addresses than IPv6
- And most routers and hosts need not change

***Keep IPv4, easier to deploy and solves more problems, including ...***

# Network filtering and traceback

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- Provides instant packet trace-ability
  - Records the route the packet takes
- Versus other schemes
  - Anti-source spoofing (ingress filtering) is not scalable
  - Statistical techniques do not respond fast enough
- Allows scalable network-based filtering
  - Push filters back along receive path to ingress points
  - Reduces flood attack to portion of bandwidth

**Research: Show WRAP/filtering can scale**

Ref. K. Argyraki, WRAP, forthcoming Ph.D. thesis

***But with source routing and FBR, there's more***

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# Instant fail-over for high availability

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- Access router maintains two or more edge-disjoint paths to destination
- Packets sent on each path
  - Recall: lots of capacity
- Duplicate suppression at receiving router
- At least one packets gets through with high probability
- Concurrent recovery of failed paths

**Research: Show FBR can scale**

Ref. D. Zhu, Feedback-based Routing, HotNets 2002, forthcoming Ph.D. thesis

# Name-based Routing

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- Route to named endpoints, not addresses
  - That's what really identifies end-system state
- Integrate naming into routing system
  - Routing system is a directory service
    - address to next hop mapping
  - Extend to provide name to next hop
- Routing protocols extended to disseminate name binding together with topology info
- Provide multi-path routing at the naming level
  - Supporting replicated sites

***True Internet routing to end-system state,  
but there's more ...***

# Highly Available Naming System

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- If you can name it, you can reach it
  - Naming in routers so no off-path dependence
- Redundancy of naming service matches  
redundancy of connectivity
  - If  $K$  multi-homed, then  $K$  separate name servers
- Attack-resistant to DDoS
- Scaling by level of indirection
  - Names to routing aggregates, routing aggregates to next-hop

**Research: Show NBR can scale**

Ref. M. Gritter, Content-based Routing, USITS  
2000, forthcoming Ph.D. thesis

# Name-based Connections

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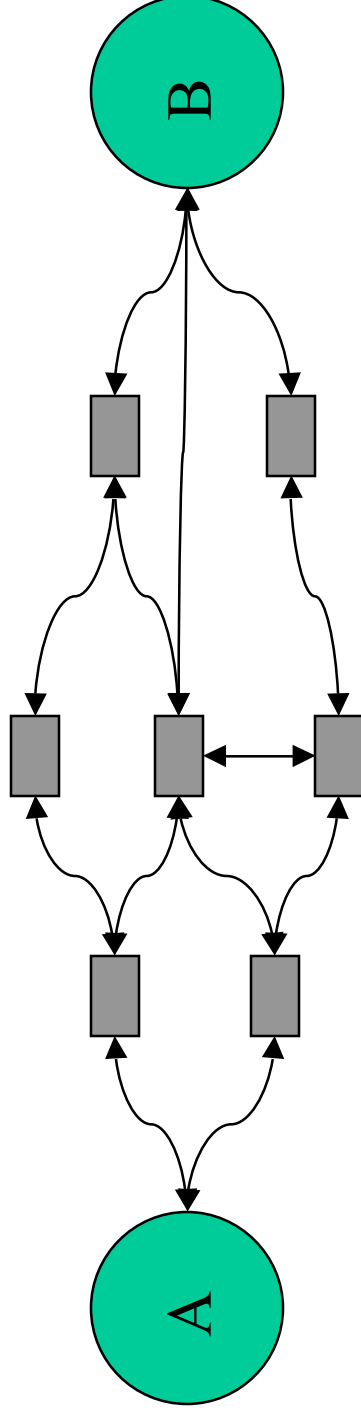
- Connection endpoint identified by name, not addr
  - i.e. specify name on connect setup and reconnect
- Name-based checksum
  - Just derive checksum base from end system names
  - Verify packet delivered to right end-system, at same cost
- Works fine with NAT
  - no dependence on addresses
  - Makes NAT state “soft”
- Deployable as a TCP option

***Provides true end-to-end reliability,***

***And allows the Internet to support NAT***

# I-ATC: Mapping application security onto Physical Security

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- True end-to-end reliability to named end systems
- Multiple disjoint redundant paths between nodes
  - Non-stop packet delivery
- Open authentication
  - Multiple messages by independent paths
  - Detection of forged attempts, like ECC
- Clear indication to network operators how to configure

***Can we trust this architecture?***

# So, adversary attacks the I-ATC

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- Crack the keys/encryption: sorry, there is none
- Forge a message:
  - Ignored because of trace-ability
  - Detected as a conflict with independent true updates
- Blow up a router:
  - no problem, use an alternative route instantly
- DDoS flooding attack:
  - repelled by network-based filtering

***Attack is, at worst, a local failure***

# Why, why, why ?

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- Why is it so hard to make architectural progress
- Named-based Internet proposed in 1991 (RFC 1287)

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# The Internet Religion

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True believers do recite:

- The Internet has been very successful so DHCP/IP/TCP/BGP/DNS must be basically right
- Minor technical extensions are the surest means to political agreement
  - DNSsec, secure BGP,
- Political solutions are solutions:
  - There are many possible technical solutions; the hard part is getting agreement, compromise is the key

***If you believe it, it will work!***

***All you need is faith***



# The Ages of the Internet Architecture

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- Age of Pioneers: 1970s
  - Bob Kahn, Vint Cerf, D. Clark, Jon Postel, Len Kleinrock, ...
  - Design and build it as “proof of concept”
- Age of Embellishers: 1980s
  - E.g. Deering/Cheriton(IP Multicast)
- Age of Religious Defenders: 1990s+
  - Return to network “transparency” – the flat earth society
  - Defending against excessive (re)invention
- Yes, we need standards, stability, etc but now it's ...

***Time for a New Age***

# The Age of Network Reason

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Architectural design based on careful specification  
on principles and properties

- Semantics
  - E.g. what does “end-to-end reliability” mean?
- Quantitative analysis of scalability

***Solid reasoning, not (just) gut instinct, faith  
and tradition***

# You may not agree entirely, but ...

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Hopefully, I've convinced you that:

- We need the right architecture and we do not have it now
  - Technical choices do matter
- We need to be faithful to the right architecture
- Many efforts are frightfully off base

So far:

- Students and I identified some of the problems
- Explored some potential solutions
  - And performed preliminary evaluation

***There is much more architectural work to do***

# Conclusions

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## The Internet architecture:

- is a success
  - Good enough to annihilate the competition
- is a disaster
  - Not good enough to handle critical systems
    - i.e. bad enough to annihilate us!

## The future Internet:

- Frightful ad hoc-ery or architecturally faithful

## The future Internet architecture

- Political sham “solutions” or science

It matters: I-ATC, You bet your life it does